

A Soft Intro to Mean Field Games

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What we are looking at

Game theory is complicated enough with only two players, but when the number explodes...

- ▶ Use 'population' arguments (congestion games...)
- ▶ Use 'algorithmic' arguments (learning...)
- ▶ Play on networks (local effects...)

Many connections between one another.

Today we are looking at way to **move to the continuous** using PDEs and mean field arguments.

If you like piña coladas...



The Beast

MFG “basic” system

$$\begin{cases} \frac{\partial u}{\partial t} + \nu \Delta u + H(x, \nabla u) + V[m] = 0 \\ \frac{\partial m}{\partial t} - \nu \Delta m + \operatorname{div} (m \nabla_p H(x, \nabla u)) = 0 \\ m|_{t=0} = m_0, u|_{t=T} = g[m_T] \end{cases}$$

Why it is worth looking at

Very general model, extends far beyond economics. The previous system admits as particular cases **famous equations** such as the Euler equations or the Vlasov ones.

Piggybacks on very active fields of research, such as **optimal transport or stochastic control**.

Has found a good deal of **numerical methods** for solving these types of games, taking it out of the theoretical forest (pedestrian crossing animation).

Derivation of the MFG system

We will derive The Beast from a simple model. Two steps:

- ▶ Agents optimize their choices given the density of other players (our m) \Rightarrow control in feedback, gives **backward** equation.
- ▶ The densities are **transported** by these choices \Rightarrow **forward** equation.

Plan

- ▶ The optimizing agent
- ▶ The transported density
- ▶ Symmetric N-player games

The agent's problem

We have a nice space Y (take \mathbb{R}^d , or call it *the beach*), with a distance d . Look at time interval $[0, T]$.

For all $t \in [0, T]$, suppose the density of the players on Y , m_t is given.

The agent controls (using $\alpha \in A$) the stochastic equation

$$dX_t = \alpha_t dt + \sqrt{2\nu} dB_t$$

where B_t is a Brownian motion in \mathbb{R}^d .

The agent's problem (cont.)

The agent solves the following **stochastic control** problem

$$\begin{cases} \inf_{\alpha} \mathbb{E} \left[\int_0^T \left(L(X_t, \alpha_t) + V[m_t](X_t) \right) dt + g[m_T](X_T) \right] \\ dX_t = \alpha_t dt + \sqrt{2\nu} dB_t \\ X_0 = x \end{cases}$$

Define Hamiltonian H

$$H(x, p) = \min_{\alpha \in A} \{ L(x, \alpha) + \alpha \cdot p \}$$

Hamilton-Jacobi-Bellman

The value function u satisfies

$$\begin{cases} \frac{\partial u}{\partial t} + \nu \Delta u + H(x, \nabla u) + V[m] = 0 \\ u|_{t=T}(x) = g[m_T](x) \end{cases}$$

A closer look at the V

V is the **cost incurred from the density of other players**. Can give a general example:

$$V[m](x) = \underbrace{V_0(x)}_{\text{I like the bar}} + \overbrace{f(x, m(x))}^{\text{I don't like people}} + \underbrace{\int_Y W(x, y)m(dy)}_{\text{But not too alone either}}$$

BUT

- ▶ With congestion term $f(x, m(x))$, no existence of equilibrium (need $V[m] \in C(Y)$)
- ▶ With distance term $\int_Y W(x, y)m(dy)$, no uniqueness (need $V[m]$ to be monotonous in m , e/g $V[m](x, m(x)) \geq V[m](y, m(y))$ if $m(x) \geq m(y)$)

Plan

- ▶ The optimizing agent
- ▶ **The transported density**
- ▶ Symmetric N-player games

Movement of agents

So far we have considered the measure m as given, but if agents move, m does too.

Agents find an **optimal control** $\bar{\alpha}(x)$ such that

$$\nabla_p H(x, p) = \bar{\alpha}(x)$$

This drift is their **motion**.

Transport of the measure

Now, for all t , $m_t = \mathcal{L}(X_t)$, with $m_0 = \mathcal{L}(X_0)$ given. So m_t is the solution of

$$\begin{cases} \int_Y \phi \, dm_t = \int_Y \mathbb{E}[\phi(X_t^x)] \, dm_0(x) \\ dX_t = \alpha_t dt + \sqrt{2\nu} dB_t \\ m_0 = \mathcal{L}(X_0) \end{cases}$$

ϕ is a test function ($\in C_c^\infty(Y)$)

Fokker-Planck equation

m_t is governed by

$$\begin{cases} \frac{\partial m}{\partial t} - \nu \Delta m + \operatorname{div}(m\bar{\alpha}) = 0 \\ m|_{t=0} = m_0 \end{cases}$$

The Beast, dissected

Clear where this system is coming from?

$$\begin{cases} \frac{\partial u}{\partial t} + \nu \Delta u + H(x, \nabla u) + V[m] = 0 \\ \frac{\partial m}{\partial t} - \nu \Delta m + \operatorname{div} (m \nabla_p H(x, \nabla u)) = 0 \\ m|_{t=0} = m_0, u|_{t=T} = g[m_T] \end{cases}$$

The linear-quadratic case

Let $L(x, \alpha) = \frac{|\alpha|^2}{2}$, then

$$H(x, p) = \min_{\alpha} \frac{|\alpha|^2}{2} + p \cdot \alpha = -\frac{|p|^2}{2}$$

The MFG system now reads

$$\begin{cases} \frac{\partial u}{\partial t} + \nu \Delta u - \frac{|\nabla u|^2}{2} + V[m] = 0 \\ \frac{\partial m}{\partial t} - \nu \Delta m - \operatorname{div}(m \nabla u) = 0 \\ m|_{t=0} = m_0, u|_{t=T} = g[m_T] \end{cases}$$

Plan

- ▶ The optimizing agent
- ▶ The transported density
- ▶ Symmetric N-player games

Symmetric N -player games

We can look at games with N players and see what happens when $N \rightarrow +\infty$.

We assume all players are **identical**, choose an action $x_i \in Y$. Cost function for player i is

$$J_i(x_i, x_{-i}) = F(x_i, x_{-i})$$

F is **symmetric** in the sense that for all permutations σ on $N - 1$, $F(x_i, x_{-i}) = F(x_i, (x_{\sigma(j)})_{j \neq i})$.

If $N \rightarrow +\infty$, then the number of variables in F also goes to infinity...

Ashes to ashes, points to measures

The trick: map positions $x = (x_1, \dots, x_N) \subset Y$ of N players to the **empirical measure** $m_x^N \in \mathcal{P}(Y)$

$$m_x^N = \frac{1}{N} \sum_{i \in N} \delta_{x_i}$$

Then if $x = y$ up to a permutation, $m_x^N = m_y^N$.

(Some like to think of x and y as belonging to the quotient space Y/\mathfrak{S}_N)

Distances on $\mathcal{P}(Y)$

Who says limits says distances, and we have a couple options for our space of measures $\mathcal{P}(Y)$.

Let m, \tilde{m} in $\mathcal{P}(Y)$. Define the **1-Wasserstein distance** by

$$W_1(m, \tilde{m}) = \begin{cases} \inf \int_{Y \times Y} d(x, y) \gamma(dx, dy) \\ \gamma \in \Pi(m, \tilde{m}) \end{cases}$$

$\Pi(m, \tilde{m})$ is the set of **couplings** between m and \tilde{m} , i/e joint distributions whose marginals are m and \tilde{m} .

Connections with **optimal transport** here...

Déblais et remblais

Problem: X and Y have the same volume, find the cheapest transportation plan between them.

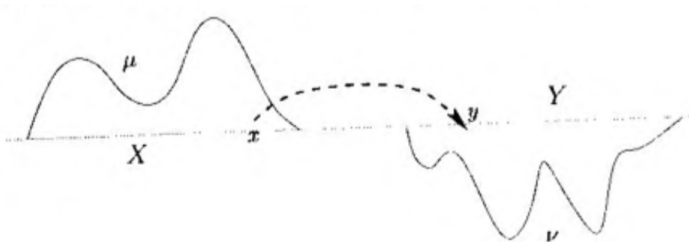


Figure: Pulled directly from *Topics on Optimal Transportation*, C. Villani

W_1 for empirical measures

To nail our transition from Y to $\mathcal{P}(Y)$, we show **our distances translate naturally**.

Take x and y in Y^N , define \bar{d} a distance such that

$$\bar{d}(x, y) = \min_{\sigma \in \mathfrak{S}_N} \frac{1}{N} \sum_{i \in N} d(x_i, y_{\sigma(i)})$$

Then $W_1(m_x^N, m_y^N) = \bar{d}(x, y)$ (uses Birkhoff's theorem on bistochastic matrices).

Looking back

- ▶ We started from a space Y of **players' positions**.
- ▶ We moved from Y to $\mathcal{P}(Y)$ by using **empirical measures**.
- ▶ We defined a distance on the space $\mathcal{P}(Y)$, the **1-Wasserstein distance**.
- ▶ This distance metrizes the **weak topology** on $\mathcal{P}(Y)$, i/e if $m_n \rightarrow^* m \Rightarrow W_1(m_n, m) \rightarrow 0$.

\Rightarrow We can take limits in $\mathcal{P}(Y)$!

The limit of functions with N variables

Let $F_N : Y^N \rightarrow \mathbb{R}$ symmetric, $F_N(x) = F_N(m_x^N)$. Under nice assumptions, there is a function $F \in C(\mathcal{P}(Y))$ such that:

$$\sup_{x \in Y^N} |F_N(m_x^N) - F(m_x^N)| \xrightarrow{N \rightarrow \infty} 0$$

Back to games

We now have a strong conceptual background to look at limits of games.

Assume each player has cost function of the form $F_i^N(x_1, \dots, x_N)$. Our previous theorem says **we can send N to infinity**.

Asymptotically, players' costs will be close to

$$F\left(x_i, \frac{1}{N-1} \sum_{j \neq i} \delta_{x_j}\right)$$

for some $F : Y \times \mathcal{P}(Y) \rightarrow \mathbb{R}$.

Nash equilibrium in mean field

We arrive at our mean field theorem.

Mean field equation

Assume that for all N , $X^N = (x_1^N, \dots, x_N^N)$ is a Nash equilibrium for the game (F_1^N, \dots, F_N^N) . Then, up to a subsequence, the sequence of empirical measures (m^N) converges to a measure \bar{m} such that

$$\int_Y F(y, \bar{m}) d\bar{m}(y) = \inf_{m \in \mathcal{P}(Y)} \int_Y F(y, m) dm(y)$$

We did the reasoning for pure strategies, can extend to mixed
 \Rightarrow **guarantees existence of NE.**

Where to go

Once we have all this, we can tie up the two parts (derivation of the MFG system and symmetric N -player games).

Assume N players, each one controls system in the form

$$\int_0^T L_i(x_1(t), \dots, x_N(t), \alpha_i(t)) dt + g_i(x_1(T), \dots, x_N(T))$$

Can assume

- ▶ $L_i(x_1(t), \dots, x_N(t), \alpha_i(t)) = \frac{|\alpha|^2}{2} + F\left(\frac{1}{N-1} \sum_{j \neq i} \delta_{x_j}\right)$
- ▶ $g_i(x_1(T), \dots, x_N(T)) = g\left(x_i, \frac{1}{N-1} \sum_{j \neq i} \delta_{x_j}\right)$

Nash equilibrium

We assume there is a map U_i^N such that

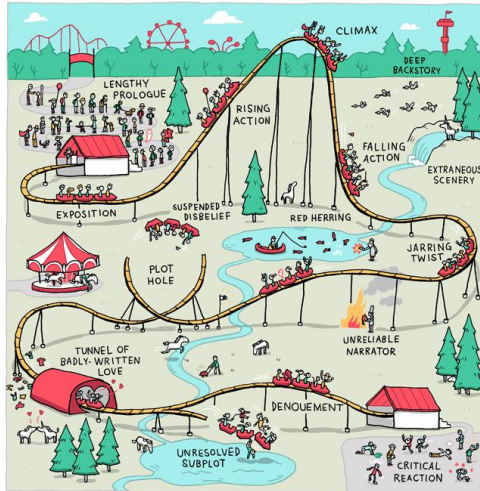
$$U_i^N(x_i, t, (x_j)_{j \neq i}) = U^N(x_i, t, (x_j)_{j \neq i})$$

and it satisfies a bunch of **HJB equations**

$$\begin{cases} \frac{\partial U_i^N}{\partial t} + \frac{1}{2} \left| \nabla_{x_i} U_i^N \right|^2 - F\left(\frac{1}{N-1} \sum_{j \neq i} \delta_{x_j}\right) + \sum_{j \neq i} \nabla_{x_j} U_j^N \cdot \nabla_{x_j} U_i^N = 0 \\ U_i^N = g_i, t = T \end{cases}$$

In other words, we have a direction for the player to follow, and the optimal control $\bar{\alpha}_i = -\nabla_{x_i} U_i^N$ is a Nash equilibrium (can prove that).

THE STORY COASTER



GRANT SNIDER

Suspended disbelief

We know we can push N to infinity and the symmetric function U^N will have a **limit** U , but we need to do it in the space of probabilities $P(Y)$.

Using estimates on $\nabla_{x_i} U^N$, we can **switch to that space** (actually using the 2-Wasserstein distance...)

We then give ourselves the initial measure m_0 and push it forward using the obtained U (and its derivatives).

Resulting system

With a bit more work we obtain:

$$\begin{cases} -\frac{\partial u}{\partial t} + \frac{1}{2}|Du|^2 = F(m) \\ \frac{\partial m}{\partial t} - \operatorname{div}(mDu) = 0 \\ m|_{t=0} = m_0, u(x, T) = g(x, m(T)) \end{cases}$$

References

Good references out there, I used:

- ▶ P.L Lions videos at the Collège de France (some are in English)
- ▶ P. Cardaliaguet's notes (online)
- ▶ Lecture notes from class taught by G. Carlier in Dauphine

Cocorico!

Thank you!